# A Tour of Proximal Algorithms 

Neal Parikh<br>Cornell University<br>March 1, 2018

## Motivation

- consider the generic convex optimization problem

$$
\begin{array}{ll}
\operatorname{minimize} & f(x) \\
\text { subject to } & x \in \mathcal{C}
\end{array}
$$

with $x \in \mathbf{R}^{n}$

- possible 'classical' approaches, depending on $f, \mathcal{C}$, and $n$ :
- unconstrained: gradient method, Newton method, BFGS
- constrained: projected gradient, primal-dual interior-point method
- potential issues:
- applies only to unconstrained problems
- assumes that objective is smooth
- requires a problem transformation that obscures problem structure
- does not scale beyond medium size problems
- does not easily support distributed data or parallel computation


## Example: $\ell_{1}$ regularization

consider the (very common) problem

$$
\operatorname{minimize} \quad f(x)+\lambda\|x\|_{1}
$$

options:

- generic subgradient method
- use a transformation like

$$
\begin{array}{ll}
\operatorname{minimize} & s+\lambda 1^{T} t \\
\text { subject to } & f(x) \leq s \\
& x_{i} \leq t_{i}, \quad i=1, \ldots, n \\
& -x_{i} \leq t_{i}, \quad i=1, \ldots, n
\end{array}
$$

and then attempt to use a symmetric cone solver

- won't work if, e.g., $f$ involves exp or log
- obscures structure, e.g., solution is only approximately sparse


## Goals

simple, general-purpose, non-heuristic methods for

- arbitrary-scale optimization (like learning/statistics with huge datasets)
- use problem structure to decompose problems into smaller/simpler pieces
want to obtain an exact, global solution to original convex problem


## Distributed SVM: Iteration 1



## Distributed SVM: Iteration 5



## Distributed SVM: Iteration 40



## This talk

interaction between three key ingredients:
(1) proximal operator of a convex function
(2) operator splitting algorithms
(3) problem transformations

## Structure and Regularization

## Outline

# $\ell_{1}$ regularization 

Examples and extensions
Proximal operators
Proximal algorithms

## Structure in variables

- often know or assume that solution to a problem is structured, e.g.,
- convex-cardinality problems
- high-dimensional statistics: assume low-dimensional structure
- prior knowledge that variables have, e.g., hierarchical or grouped structure
- handle by solving a problem with two conceptual components:
- main objective of interest (model fit, satisfying constraints, ...)
- regularization term that encourages assumed form of structure
- possible structure of interest includes sparsity, low rank, ...


## this talk:

(1) selecting regularization to promote assumed structure
(2) many examples and applications (i.e., sparsify everything in sight)
(3) solving the resulting optimization problems

## Geometric interpretation



get sparsity/structure when corners/kinks appear at sparse/structured points e.g., quadratic cone, linear functions on prob. simplex, nuclear norm, ...

## Convex envelope interpretation

- convex envelope of (nonconvex) $f$ is the largest convex underestimator $g$
- i.e., the best convex lower bound to a function

- example: $\ell_{1}$ is the envelope of card (on unit $\ell_{\infty}$ ball)
- example: $\|\cdot\|_{*}$ is the envelope of rank (on unit spectral norm ball)
- various characterizations: e.g., $f^{* *}$ or convex hull of epigraph


## Penalty function interpretation

- compared to ridge penalty $\|\cdot\|_{2}^{2}$, using $\ell_{1}$ does two things:
(1) higher emphasis on small values to go to exactly zero
(2) lower emphasis on avoiding very large values
- thus useful for obtaining sparse or robust solutions to problems


## Atomic norm interpretation

## (Chandrasekaran, Recht, Parrilo, Willsky)

- convex surrogates for measures of 'simplicity'
- suppose underlying parameter vector or signal $x \in \mathbf{R}^{n}$ given by

$$
x=\sum_{i=1}^{k} c_{i} a_{i}, \quad a_{i} \in \mathcal{A}, c_{i} \geq 0
$$

where $\mathcal{A}$ is set of 'atoms' and $k \ll n$ (d.f. $\ll$ ambient dimension)

- if $\mathcal{A}$ is usual basis vectors, model says that $x$ is $k$-sparse, and

$$
\operatorname{conv}(\mathcal{A})=\text { unit } \ell_{1} \text { ball }
$$

- then, e.g., minimize $\|x\|_{1}$ subject to $y=F x$


## Outline

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## Sparse design

- find sparse design vector $x$ satisfying specifications

$$
\begin{array}{ll}
\operatorname{minimize} & \|x\|_{1} \\
\text { subject to } & x \in \mathcal{C}
\end{array}
$$

- zero values of $x$ simplify design or correspond to unneeded components
- when $\mathcal{C}=\{x \mid A x=b\}$, called basis pursuit or sparse coding
- e.g., find sparse representation of signal $b$ in 'dictionary' or 'overcomplete basis' given by columns of $A$


## Sparse regression

- fit $b \in \mathbf{R}^{m}$ as linear combination of a subset of regressors

$$
\operatorname{minimize} \quad(1 / 2)\|A x-b\|_{2}^{2}+\lambda\|x\|_{1}
$$

- zero values of $x$ indicate features not predictive of the response
- also known as the lasso
- easily generalizes to other losses (e.g., sparse logistic regression)


## Sparse regression


original

quadratic regularization

## Sparse regression


original

$\ell_{1}$ norm

## Estimation with outliers

- measurements $y_{i}=a_{i}^{T} x+v_{i}+w_{i}$
- $v_{i}$ is Gaussian noise (small), $w$ is a sparse outlier vector (big)
- if $\mathcal{O}=\left\{i \mid w_{i} \neq 0\right\}$ is set of outliers, MLE given by

$$
\begin{array}{ll}
\operatorname{minimize} & \sum_{i \notin \mathcal{O}}\left(y_{i}-a_{i}^{T} x\right)^{2} \\
\text { subject to } & |\mathcal{O}| \leq k
\end{array}
$$

- convex approximation given by

$$
\operatorname{minimize} \quad(1 / 2)\|y-A x-w\|_{2}^{2}+\lambda\|w\|_{1}
$$

- same idea used in support vector machine


## Linear classifier with fewest errors

- want linear classifier $b \approx \operatorname{sign}\left(a^{T} x+s\right)$ from $\left(a_{i}, b_{i}\right) \in \mathbf{R}^{n} \times\{-1,1\}$
- error corresponds to negative margin: $b_{i}\left(a_{i}^{T} x+s\right) \leq 0$
- find $x, s$ that give fewest classification errors:

$$
\begin{array}{ll}
\operatorname{minimize} & \|t\|_{1} \\
\text { subject to } & b_{i}\left(a_{i}^{T} x+s\right)+t_{i} \geq 1, \quad i=1, \ldots m
\end{array}
$$

with variables $x, s, t$

- close to a support vector machine
- can generalize to other convex feasibility problems


## Elastic net

## (Zou \& Hastie)

- problem:

$$
\operatorname{minimize} \quad f(x)+\lambda\|x\|_{1}+(1-\lambda)\|x\|_{2}^{2}
$$

i.e., use both ridge and lasso penalties

- attempts to overcome the following potential drawbacks of the lasso:
- lasso selects at most (\# examples) variables
- given group of very correlated features, lasso often picks one arbitrarily
- here, strongly correlated predictors are jointly included or not
- (in practice, need to do some rescaling above)


## Fused lasso

## (Tibshirani et al.; Rudin, Osher, Fatemi)

- problem:

$$
\operatorname{minimize} \quad f(x)+\lambda_{1}\|x\|_{1}+\lambda_{2} \sum_{j=2}^{n}\left|x_{j}-x_{j-1}\right|
$$

i.e., encourage $x$ to be both sparse and piecewise constant

- special case: total variation denoising (set $\lambda_{1}=0$ )
- used in biology (e.g., gene expression) and signal reconstruction
- can also write penalty as $\|D x\|_{1}$; could consider other matrices


## Total variation denoising




120 linear measurements and $31 \times 31=961$ variables (' $8 \times$ undersampled')

## Total variation denoising




120 linear measurements and $31 \times 31=961$ variables (' $8 \times$ undersampled')

## Group lasso

(e.g., Yuan \& Lin; Meier, van de Geer, Bühlmann; Jacob, Obozinski, Vert)

- problem:
minimize $\quad f(x)+\lambda \sum_{i=1}^{N}\left\|x_{i}\right\|_{2}$
i.e., like lasso, but require groups of variables to be zero or not
- also called $\ell_{1,2}$ mixed norm regularization
- related to multiple kernel learning via duality (see Bach et al.)


## Joint covariate selection for multi-task learning

 (Obozinski, Taskar, Jordan)- want to fit parameters $x^{k} \in \mathbf{R}^{p}$ for each of multiple datasets $\mathcal{D}^{k}$
- either use feature $j$ in all tasks or none of them
- let $x_{j}=\left(x_{j}^{1}, \ldots, x_{j}^{K}\right)$ for $j=1, \ldots, p$
- problem:

$$
\operatorname{minimize} \quad \sum_{k=1}^{K} f^{k}\left(x^{k}\right)+\lambda \sum_{j=1}^{p}\left\|x_{j}\right\|_{2}
$$

with variables $x^{1}, \ldots, x^{K} \in \mathbf{R}^{p}$

## Structured group lasso

(Jacob, Obozinski, Vert; Bach et al.; Zhao, Rocha, Yu; ...)

- problem:

$$
\operatorname{minimize} \quad f(x)+\sum_{i=1}^{N} \lambda_{i}\left\|x_{g_{i}}\right\|_{2}
$$

where $g_{i} \subseteq[n]$ and $\mathcal{G}=\left\{g_{1}, \ldots, g_{N}\right\}$

- like group lasso, but the groups can overlap arbitrarily
- particular choices of groups can impose 'structured' sparsity
- e.g., topic models, selecting interaction terms for (graphical) models, tree structure of gene networks, fMRI data
- generalizes to the composite absolute penalties family:

$$
r(x)=\left\|\left(\left\|x_{g_{1}}\right\|_{p_{1}}, \ldots,\left\|x_{g_{N}}\right\|_{p_{N}}\right)\right\|_{p_{0}}
$$

## Structured group lasso

(Jacob, Obozinski, Vert; Bach et al.; Zhao, Rocha, Yu; ...)

contiguous selection:


- $\mathcal{G}=\{\{1\},\{5\},\{1,2\},\{4,5\},\{1,2,3\},\{3,4,5\},\{1,2,3,4\},\{2,3,4,5\}\}$
- nonzero variables are contiguous in $x$, e.g., $x^{\star}=(0, *, *, 0,0)$
- can extend the same idea to higher dimensions (e.g., select rectangles)
- e.g., time series, tumor diagnosis, ...


## Structured group lasso

(Jacob, Obozinski, Vert; Bach et al.; Zhao, Rocha, Yu; ...)
hierarchical selection:


- $\mathcal{G}=\{\{4\},\{5\},\{6\},\{2,4\},\{3,5,6\},\{1,2,3,4,5,6\}\}$
- nonzero variables form a rooted and connected subtree
- if node is selected, so are its ancestors
- if node is not selected, neither are its descendants


## Matrix decomposition

- problem:

$$
\begin{array}{ll}
\operatorname{minimize} & f_{1}\left(X_{1}\right)+\cdots+f_{N}\left(X_{N}\right) \\
\text { subject to } & X_{1}+\cdots+X_{N}=A
\end{array}
$$

- many choices for the $f_{i}$ :
- squared Frobenius norm (least squares)
- entrywise $\ell_{1}$ norm (sparse matrix)
- nuclear norm (low rank)
- sum-\{row,column\}-norm (group lasso)
- elementwise constraints (fixed sparsity pattern, nonnegative, ...)
- semidefinite cone constraint


# Low rank matrix completion 

(Candès \& Recht; Recht, Fazel, Parrilo)

- problem:

$$
\begin{array}{ll}
\operatorname{minimize} & \|X\|_{*} \\
\text { subject to } & X_{i j}=A_{i j}, \quad(i, j) \in \mathcal{D}
\end{array}
$$

i.e., find low rank matrix that agrees with observed entries

- e.g., Netflix problem


## Robust PCA

## (Candès et al.; Chandrasekaran et al.)

- regular PCA is the (nonconvex but solvable) problem

$$
\begin{array}{ll}
\operatorname{minimize} & \|A-L\|_{2} \\
\text { subject to } & \operatorname{rank}(L) \leq k
\end{array}
$$

i.e., recover rank $k$ matrix $L_{0}$ if $A=L_{0}+N_{0}$, where $N_{0}$ is noise

- if matrix also has some sparse but large noise, instead solve

$$
\begin{array}{ll}
\operatorname{minimize} & \|L\|_{*}+\lambda\|S\|_{1} \\
\text { subject to } & L+S=A
\end{array}
$$

i.e., recover low rank $L$ and sparse corruption $S$ if $A=L_{0}+S_{0}+N_{0}$

- sparse + low rank decomposition has other applications (e.g., vision, video segmentation, background subtraction, biology, indexing)


## Robust PCA

(Candès et al.; Chandrasekaran et al.)


## Outline

## $\ell_{1}$ regularization

## Examples and extensions

Proximal operators
Proximal algorithms

## Proximal operator

(Martinet; Moreau; Rockafellar)

- proximal operator of $f: \mathbf{R}^{n} \rightarrow \mathbf{R} \cup\{+\infty\}$ is

$$
\operatorname{prox}_{\lambda f}(v)=\underset{x}{\operatorname{argmin}}\left(f(x)+(1 / 2 \lambda)\|x-v\|_{2}^{2}\right)
$$

with parameter $\lambda>0$

- $f$ may be nonsmooth, have embedded constraints, ...
- can evaluate with standard methods like BFGS, but often has an analytical solution or simple specialized linear-time algorithm
- many interpretations
- example: proximal operator of $I_{\mathcal{C}}$ is $\Pi_{\mathcal{C}}$ (generalized projection)


## Polyhedra

- projection onto polyhedron $\mathcal{C}=\{x \mid A x=b, C x \leq d\}$ is a QP
- projection onto affine set $\mathcal{C}=\{x \mid A x=b\}$ is a linear operator
- box or hyperrectangle $\mathcal{C}=\{x \mid l \preceq x \preceq u\}$ :

$$
\left(\Pi_{\mathcal{C}}(v)\right)_{k}= \begin{cases}l_{k} & v_{k} \leq l_{k} \\ v_{k} & l_{k} \leq v_{k} \leq u_{k} \\ u_{k} & v_{k} \geq u_{k}\end{cases}
$$

- also simple methods for hyperplanes, halfspaces, simplexes, ...


## Quadratic functions

- if $f(x)=(1 / 2) x^{T} P x+q^{T} x+r$, then

$$
\operatorname{prox}_{\lambda f}(v)=(I+\lambda P)^{-1}(v-\lambda q)
$$

- if evaluating repeatedly with different arguments $v$ :
- dense direct method: $O\left(n^{3}\right)$ flops first time and then $O\left(n^{2}\right)$
- iterative method (CG, LSQR, ...): warm start beginning at $v$


## Moreau envelope

- Moreau envelope or Moreau-Yosida regularization of $f$ is

$$
M_{\lambda f}(v)=\inf _{x}\left(f(x)+(1 / 2 \lambda)\|x-v\|_{2}^{2}\right)
$$

- a smoothed or regularized form of $f$ :
- always has full domain
- always continuously differentiable
- has the same minimizers as $f$
- proximal operator is gradient step for Moreau envelope:

$$
\operatorname{prox}_{\lambda f}(x)=x-\lambda \nabla M_{\lambda f}(x)
$$

## Moreau envelope

- motivation: in general, $\varphi^{*}$ is smooth when $\varphi$ is strongly convex
- can show that

$$
M_{f}=\left(f^{*}+(1 / 2)\|\cdot\|_{2}^{2}\right)^{*}
$$

so Moreau envelope obtains a smooth approximation via
(1) taking conjugate
(2) regularizing to get a strongly convex function
(3) taking conjugate again

- example: Moreau envelope of $|\cdot|$ is the Huber loss function

Moreau envelope: Huber loss


## Moreau decomposition

- following relation always holds:

$$
v=\operatorname{prox}_{f}(v)+\operatorname{prox}_{f *}(v)
$$

- main link between proximal operators and duality
- a generalization of orthogonal decomposition induced by subspace $L$ :

$$
v=\Pi_{L}(v)+\Pi_{L^{\perp}}(v)
$$

follows from Moreau decomposition and $\left(I_{L}\right)^{*}=I_{L^{\perp}}$

## Norms and norm balls

- in general: if $f=\|\cdot\|$ and $\mathcal{B}$ is unit ball of dual norm, then

$$
\operatorname{prox}_{\lambda f}(v)=v-\lambda \Pi_{\mathcal{B}}(v / \lambda)
$$

- if $f=\|\cdot\|_{2}$ and $\mathcal{B}$ is the unit $\ell_{2}$ ball, then

$$
\begin{aligned}
\Pi_{\mathcal{B}}(v) & = \begin{cases}v /\|v\|_{2} & \|v\|_{2}>1 \\
v & \|v\|_{2} \leq 1\end{cases} \\
\operatorname{prox}_{\lambda f}(v) & = \begin{cases}\left(1-\lambda /\|v\|_{2}\right) v & \|v\|_{2} \geq \lambda \\
0 & \|v\|_{2}<\lambda\end{cases}
\end{aligned}
$$

sometimes called 'block soft thresholding' operator

## Norms and norm balls

- if $f=\|\cdot\|_{1}$ and $\mathcal{B}$ is the unit $\ell_{\infty}$ ball, then

$$
\left(\Pi_{\mathcal{B}}(v)\right)_{i}= \begin{cases}1 & v_{i}>1 \\ v_{i} & \left|v_{i}\right| \leq 1 \\ -1 & v_{i}<-1\end{cases}
$$

lets us derive (elementwise) soft thresholding

$$
\operatorname{prox}_{\lambda f}(v)=(v-\lambda)_{+}-(-v-\lambda)_{+}= \begin{cases}v_{i}-\lambda & v_{i} \geq \lambda \\ 0 & \left|v_{i}\right| \leq \lambda \\ v_{i}+\lambda & v_{i} \leq-\lambda\end{cases}
$$

- if $f=\|\cdot\|_{\infty}$ and $\mathcal{B}$ is unit $\ell_{1}$ ball, simple algorithms available


## Soft thresholding



## Matrix functions

- suppose convex $F: \mathbf{R}^{m \times n} \rightarrow \mathbf{R}$ is orthogonally invariant:

$$
F(Q X \tilde{Q})=F(X)
$$

for all orthogonal $Q, \tilde{Q}$

- then $F=f \circ \sigma$ and

$$
\operatorname{prox}_{\lambda F}(A)=U \operatorname{diag}\left(\operatorname{prox}_{\lambda f}(d)\right) V^{T}
$$

where $A=U \operatorname{diag}(d) V^{T}$ is the SVD of $A$ and $\sigma(A)=d$

- e.g., $F=\|\cdot\|_{*}$ has $f=\|\cdot\|_{1}$ so $\operatorname{prox}_{\lambda F}$ is 'singular value thresholding'


## Outline

## $\ell_{1}$ regularization

Examples and extensions
Proximal operators
Proximal algorithms

## Proximal gradient method

(e.g., Levitin \& Polyak; Mercier; Chen \& Rockafellar; Combettes; Tseng)

- problem form

$$
\operatorname{minimize} \quad f(x)+g(x)
$$

where $f$ is smooth and $g: \mathbf{R}^{n} \rightarrow \mathbf{R} \cup\{+\infty\}$ is closed proper convex

- method:

$$
x^{k+1}:=\operatorname{prox}_{\lambda^{k} g}\left(x^{k}-\lambda^{k} \nabla f\left(x^{k}\right)\right)
$$

- special case: projected gradient method (take $g=I_{\mathcal{C}}$ )


## Accelerated proximal gradient method

## (Nesterov; Beck \& Teboulle; Tseng)

- problem form

$$
\operatorname{minimize} \quad f(x)+g(x)
$$

where $f$ is smooth and $g: \mathbf{R}^{n} \rightarrow \mathbf{R} \cup\{+\infty\}$ is closed proper convex

- method:

$$
\begin{aligned}
y^{k+1} & :=x^{k}+\omega^{k}\left(x^{k}-x^{k-1}\right) \\
x^{k+1} & :=\operatorname{prox}_{\lambda^{k} g}\left(y^{k+1}-\lambda^{k} \nabla f\left(y^{k+1}\right)\right)
\end{aligned}
$$

works for, e.g., $\omega^{k}=k /(k+3)$ and particular $\lambda^{k}$

- faster in both theory and practice


## ADMM

## (e.g., Gabay \& Mercier; Glowinski \& Marrocco; Boyd et al.)

- problem form

$$
\operatorname{minimize} \quad f(x)+g(x)
$$

where $f, g: \mathbf{R}^{n} \rightarrow \mathbf{R} \cup\{+\infty\}$ are closed proper convex

- method:

$$
\begin{aligned}
x^{k+1} & :=\operatorname{prox}_{\lambda f}\left(z^{k}-u^{k}\right) \\
z^{k+1} & :=\operatorname{prox}_{\lambda g}\left(x^{k+1}+u^{k}\right) \\
u^{k+1} & :=u^{k}+x^{k+1}-z^{k+1}
\end{aligned}
$$

- basically, always works


## Examples

- (accelerated) proximal gradient for elastic net:
(1) gradient step for smooth loss (e.g., logistic, least squares, ...)
(2) shrinkage and elementwise soft thresholding
- ADMM for multi-task learning with joint covariate selection:
(1) evaluate prox $_{f^{k}}$ (in parallel for each dataset)
(2) block soft thresholding (in parallel for each feature)
(3) dual update
- ADMM for robust PCA:
(1) singular value thresholding
(2) elementwise soft thresholding
(3) dual update


## Distributed Optimization and Statistical Learning

## Goals

simple and robust methods for

- arbitrary-scale optimization
- machine learning/statistics with huge datasets
- decentralized optimization
- have devices/agents coordinate to solve problems by message passing


## Distributed SVM: Iteration 1



## Distributed SVM: Iteration 5



## Distributed SVM: Iteration 40



## Outline

## Operator splitting

Applications
Block splitting
Conclusions

## Operator splitting

- the most useful proximal methods use the idea of operator splitting
- these algorithms minimize $f+g$ only using prox $_{f}$ and/or prox $_{g}$
- useful when $f$ and $g$ each have useful structure separately
- very common in statistical applications: loss + regularizer
- transform problem (if needed) so when an operator splitting method is applied, breaks apart into small pieces with simple proximal operators


## Separable sum

- if $f$ is block separable, so $f(x)=\sum_{i=1}^{N} f_{i}\left(x_{i}\right)$, then

$$
\left(\operatorname{prox}_{f}(v)\right)_{i}=\operatorname{prox}_{f_{i}}\left(v_{i}\right), \quad i=1, \ldots, N
$$

- key to parallel/distributed proximal algorithms
- for $f=\|\cdot\|_{1}$, get soft thresholding

$$
\operatorname{prox}_{\lambda f}(v)=(v-\lambda)_{+}-(-v-\lambda)_{+}= \begin{cases}v_{i}-\lambda & v_{i} \geq \lambda \\ 0 & \left|v_{i}\right| \leq \lambda \\ v_{i}+\lambda & v_{i} \leq-\lambda\end{cases}
$$

## ADMM

(Douglas-Rachford 55, Gabay-Mercier 76, Glowinski-Marrocco 76)

$$
\operatorname{minimize} \quad f(x)+g(x)
$$

$f, g: \mathbf{R}^{n} \rightarrow \mathbf{R} \cup\{+\infty\}$ are closed proper convex

- method:

$$
\begin{aligned}
x^{k+1} & :=\operatorname{prox}_{\lambda f}\left(z^{k}-u^{k}\right) \\
z^{k+1} & :=\operatorname{prox}_{\lambda g}\left(x^{k+1}+u^{k}\right) \\
u^{k+1} & :=u^{k}+x^{k+1}-z^{k+1}
\end{aligned}
$$

- always converges (if problem is solvable)


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## Convex feasibility

- problem:

$$
\text { find } \quad x \in \mathcal{C} \cap \mathcal{D}
$$

- rewrite as

$$
\operatorname{minimize} \quad I_{\mathcal{C}}(x)+I_{\mathcal{D}}(x)
$$

- ADMM:

$$
\begin{aligned}
x^{k+1} & :=\Pi_{\mathcal{C}}\left(z^{k}-u^{k}\right) \\
z^{k+1} & :=\Pi_{\mathcal{D}}\left(x^{k+1}+u^{k}\right) \\
u^{k+1} & :=u^{k}+x^{k+1}-z^{k+1}
\end{aligned}
$$

## Positive semidefinite matrix completion

- given $A \in \mathbf{S}^{n}$ with $(i, j) \in \mathcal{K}$ known, fill in missing entries so in $\mathbf{S}_{+}^{n}$
- splitting:

$$
\mathcal{C}=\mathbf{S}_{+}^{n}, \quad \mathcal{D}=\left\{X \mid X_{i j}=A_{i j},(i, j) \in \mathcal{K}\right\}
$$

- projection onto $\mathcal{C}$ : find eigendecomposition $X=\sum_{i} \alpha_{i} v_{i} v_{i}^{T}$, then

$$
\Pi_{\mathcal{C}}=\sum_{i=1}^{n} \max \left\{0, \alpha_{i}\right\} v_{i} v_{i}^{T}
$$

- projection onto $\mathcal{D}$ sets entries to known values


## Positive semidefinite matrix completion



- example with $50 \times 50$ matrix missing half its entries
- blue: alternating projections; red: ADMM
- $X^{k} \in \mathcal{C}, Z^{k} \in \mathcal{D}$


## Lasso

- to minimize $(1 / 2)\|A x-b\|_{2}^{2}+\gamma\|x\|_{1}$ :

$$
\begin{aligned}
x^{k+1} & :=\left(I+\lambda A^{T} A\right)^{-1}\left(z^{k}-u^{k}-\lambda A^{T} b\right) \\
z^{k+1} & :=\operatorname{prox}_{\lambda \gamma\|\cdot\|_{1}}\left(x^{k+1}+u^{k}\right) \\
u^{k+1} & :=u^{k}+x^{k+1}-z^{k+1}
\end{aligned}
$$

- faster implementations:
- matrix inversion lemma
- factorization caching
- warm start
- adjusting $\lambda$
- easily generalizes, e.g., sparse inverse covariance selection


## Global consensus optimization

- minimize $f(x)=\sum_{i=1}^{N} f_{i}(x)$; e.g., $f_{i}$ is loss for $i$ th shard of data
- handle each shard separately via splitting

$$
\operatorname{minimize} \quad \sum_{i=1}^{N} f_{i}\left(x_{i}\right)+I_{\mathcal{C}}\left(x_{1}, \ldots, x_{N}\right)
$$

with consensus set $\mathcal{C}=\left\{\left(x_{1}, \ldots, x_{N}\right) \in \mathbf{R}^{n N} \mid x_{1}=x_{2}=\cdots=x_{N}\right\}$

- ADMM simplifies to

$$
\begin{aligned}
x_{i}^{k+1} & :=\operatorname{prox}_{f_{i}}\left(\bar{x}^{k}-u_{i}^{k}\right) \\
u_{i}^{k+1} & :=u_{i}^{k}+x_{i}^{k+1}-\bar{x}^{k+1}
\end{aligned}
$$

- intuition: $u_{i}$ measures deviation from average, proximal penalty balances minimizing $f_{i}$ while pulling towards average


## Distributed lasso via consensus optimization

$$
\begin{aligned}
x_{i}^{k+1} & :=\left(I+\lambda A_{i}^{T} A_{i}\right)^{-1}\left(z^{k}-u_{i}^{k}-\lambda A_{i}^{T} b\right) \\
z^{k+1} & :=\operatorname{prox}_{(\lambda \gamma / N)\|\cdot\|_{1}}\left(\bar{x}^{k+1}+\bar{u}^{k}\right) \\
u_{i}^{k+1} & :=u_{i}^{k}+x_{i}^{k+1}-z^{k+1}
\end{aligned}
$$

## Distributed global consensus optimization with MPI

initialize $N$ processes, along with $x_{i}, u_{i}, z$. repeat until converged

1. Update $u_{i}:=u_{i}+x_{i}-z$.
2. Update $x_{i}:=\operatorname{prox}_{\lambda f_{i}}\left(z-u_{i}\right)$.
3. Let $w:=x_{i}+u_{i}$.
4. Allreduce $w$.
5. Update $z:=\operatorname{prox}_{(\lambda / N) g}(w / N)$.
(SPMD: code runs on each separate machine, so $i$ refers to 'local' version)

## Distributed lasso example

- example with dense $A \in \mathbf{R}^{400000 \times 8000}$ (roughly 30 GB of data)
- distributed solver written in C using MPI and GSL
- no optimization or tuned libraries
- split into 80 subsystems across 10 (8-core) machines on Amazon EC2
- computation times
loading data 30s
factorization 5 m
subsequent ADMM iterations 0.5-2s
lasso solve (about 15 ADMM iterations) 5-6m


## Matrix decomposition

- decompose matrix $A$ into sum of 'simple' components:

$$
\begin{array}{ll}
\operatorname{minimize} & f_{1}\left(X_{1}\right)+f_{2}\left(X_{2}\right)+\cdots+f_{N}\left(X_{N}\right) \\
\text { subject to } & A=X_{1}+X_{2}+\cdots+X_{N}
\end{array}
$$

- penalty functions can include
- squared Frobenius norm
- entrywise $\ell_{1}$ norm
- sum-\{row,column\} norm
- indicator of elementwise constraints
- indicator of semidefinite cone
- nuclear norm


## Matrix decomposition via ADMM

- splitting:

$$
\operatorname{minimize} \quad \sum_{i=1}^{N} f_{i}\left(X_{i}\right)+I_{\mathcal{C}}\left(X_{1}, \ldots, X_{N}\right)
$$

with equilibrium set $\mathcal{C}=\left\{\left(X_{1}, \ldots, X_{N}\right) \mid A=X_{1}+\cdots+X_{N}\right\}$

- ADMM simplifies to:

$$
\begin{aligned}
X_{i}^{k+1} & :=\operatorname{prox}_{\lambda f_{i}}\left(X_{i}^{k}-\bar{X}^{k}+(1 / N) A-U^{k}\right) \\
U^{k+1} & :=U^{k}+\bar{X}^{k+1}-(1 / N) A
\end{aligned}
$$

## Matrix decomposition results

problem: decompose $A=$ rank $4+$ sparse + small Gaussian noise

| Method | $m$ | $n$ | Iterations | Time (s) |
| :--- | :--- | :--- | :--- | :--- |
| CVX | 10 | 30 | 15 | 1.11 |
| ADMM | 10 | 30 | 45 | 0.02 |
| CVX | 20 | 50 | 17 | 2.54 |
| ADMM | 20 | 50 | 42 | 0.03 |
| CVX | 40 | 80 | 20 | 108.14 |
| ADMM | 40 | 80 | 36 | 0.07 |
| ADMM | 100 | 200 | 38 | 0.58 |
| ADMM | 500 | 1000 | 42 | 35.56 |

note: last instance has 1.5 M variables and 500 K constraints

## Outline

## Operator splitting

Applications
Block splitting
Conclusions

## Graph form problems

- graph form problem:

$$
\begin{array}{ll}
\operatorname{minimize} & f(y)+g(x) \\
\text { subject to } & y=A x
\end{array}
$$

where $A \in \mathbf{R}^{m \times n}, f: \mathbf{R}^{m} \rightarrow \mathbf{R} \cup\{+\infty\}, g: \mathbf{R}^{n} \rightarrow \mathbf{R} \cup\{+\infty\}$

- $x$ and $y$ must lie in the graph $\left\{(x, y) \in \mathbf{R}^{m+n} \mid y=A x\right\}$ of $A$
- refer to $x$ as 'inputs' and $y$ as 'outputs'
- $f$ and $g$ can encode constraints


## Example: cone programming

- cone program in standard form:

$$
\begin{array}{ll}
\operatorname{minimize} & c^{T} x \\
\text { subject to } & A x=b \\
& x \in \mathcal{K}
\end{array}
$$

where $\mathcal{K}$ is a convex cone

- in graph form, let

$$
f(y)=I_{\{b\}}(y), \quad g(x)=c^{T} x+I_{\mathcal{K}}(x)
$$

where $I_{\mathcal{C}}$ is the indicator function of the convex set $\mathcal{C}$

- e.g., symmetric cone program when $\mathcal{K}$ is a product of $\mathbf{R}_{+}^{n}, \mathbf{Q}^{n}, \mathbf{S}_{+}^{n}$


## Example: loss minimization

- many statistics/ML problems take the form

$$
\operatorname{minimize} \quad l(A x-b)+r(x)
$$

- in graph form, let

$$
f(y)=l(y-b), \quad g(x)=r(x)
$$

- e.g., obtain the lasso with $l(u)=(1 / 2)\|u\|_{2}^{2}$ and $r(v)=\gamma\|v\|_{1}$
- can similarly express linear SVM, MLE/MAP in exponential families, ...


## ADMM

- consider generic constrained convex program

$$
\begin{array}{ll}
\operatorname{minimize} & f(z) \\
\text { subject to } & z \in \mathcal{C}
\end{array}
$$

- ADMM:

$$
\begin{array}{lll}
z^{k+1 / 2} & :=\operatorname{prox}_{f}\left(z^{k}-\tilde{z}^{k}\right) & \text { // prox } \\
z^{k+1} & :=\Pi_{\mathcal{C}}\left(z^{k+1 / 2}+\tilde{z}^{k}\right) & \text { // projection } \\
\tilde{z}^{k+1} & :=\tilde{z}^{k}+z^{k+1 / 2}-z^{k+1} & \text { // dual update }
\end{array}
$$

- converges under very general conditions


## Graph projection splitting

- applying this form of ADMM to graph form problem gives

$$
\begin{aligned}
x^{k+1 / 2} & :=\operatorname{prox}_{g}\left(x^{k}-\tilde{x}^{k}\right) \\
y^{k+1 / 2} & :=\operatorname{prox}_{f}\left(y^{k}-\tilde{y}^{k}\right) \\
\left(x^{k+1}, y^{k+1}\right) & :=\Pi_{A}\left(x^{k+1 / 2}+\tilde{x}^{k}, y^{k+1 / 2}+\tilde{y}^{k}\right) \\
\tilde{x}^{k+1} & :=\tilde{x}^{k}+x^{k+1 / 2}-x^{k+1} \\
\tilde{y}^{k+1} & :=\tilde{y}^{k}+y^{k+1 / 2}-y^{k+1}
\end{aligned}
$$

- $\Pi_{A}$ is called graph projection and denotes projection onto graph of $A$
- important: $f$ and $g$ never interact directly with $A$, i.e., $\Pi_{A}$ is the only operation that touches the data


## Graph projection

- evaluating $\Pi_{A}(c, d)$ involves solving

$$
\begin{array}{ll}
\operatorname{minimize} & (1 / 2)\|x-c\|_{2}^{2}+(1 / 2)\|y-d\|_{2}^{2} \\
\text { subject to } & y=A x
\end{array}
$$

- reduce to solving (quasidefinite) KKT system

$$
\left[\begin{array}{cc}
I & A^{T} \\
A & -I
\end{array}\right]\left[\begin{array}{l}
x \\
y
\end{array}\right]=\left[\begin{array}{c}
c+A^{T} d \\
0
\end{array}\right]
$$

## Implementing graph projections

- eliminate $x$, then solve for $y$

$$
\begin{aligned}
& y:=\left(I+A A^{T}\right)^{-1}\left(A c+A A^{T} d\right) \\
& x:=c+A^{T}(d-y)
\end{aligned}
$$

- eliminate $y$, then solve for $x$

$$
\begin{aligned}
& x:=\left(I+A^{T} A\right)^{-1}\left(c+A^{T} d\right) \\
& y:=A x
\end{aligned}
$$

- when $A$ is dense, prefer first when $A$ is fat and second when $A$ is skinny
- factor the relevant coefficient matrix with, e.g., Cholesky factorization
- see paper for more on other situations (e.g., sparse $A$ )
- key point: does not depend on $\lambda, f, g$


## Example

- graph projection splitting algorithm for the lasso:

$$
\begin{aligned}
x^{k+1 / 2} & :=\operatorname{prox}_{\lambda \gamma\|\cdot\|_{1}}\left(x^{k}-\tilde{x}^{k}\right) \\
y^{k+1 / 2} & :=(1 /(1+\lambda))\left(y^{k}-\tilde{y}^{k}\right) \\
\left(x^{k+1}, y^{k+1}\right) & :=\Pi_{A}\left(x^{k+1 / 2}+\tilde{x}^{k}, y^{k+1 / 2}+\tilde{y}^{k}\right) \\
\tilde{x}^{k+1} & :=\tilde{x}^{k}+x^{k+1 / 2}-x^{k+1} \\
\tilde{y}^{k+1} & :=\tilde{y}^{k}+y^{k+1 / 2}-y^{k+1}
\end{aligned}
$$

- here, all operations are trivial except for $\Pi_{A}$
- cache factorization needed to evaluate $\Pi_{A}$ and reuse after first iteration
- regularization path: reuse across solves, varying $\gamma$
- model comparison: reuse across solves, varying $f$ or $g$


## Sample implementation: lasso

```
prox_f = @(v,lambda) (1/(1 + lambda))*(v - b) + b;
prox_g = @(v,lambda) (max(0, v - lambda) - max(0, -v - lambda));
AA = A*A';
L = chol(eye(m) + AA);
for iter = 1:MAX_ITER
    xx = prox_g(xz - xt, lambda);
    yx = prox_f(yz - yt, lambda);
    yz = L \ (L' \ (A*(xx + xt) + AA*(yx + yt)));
    xz = xx + xt + A'*(yx + yt - yz);
    xt = xt + xx - xz;
    yt = yt + yx - yz;
end
```


## Numerical example: regularization path for lasso

- dense $A \in \mathbf{R}^{5000 \times 8000}$
- 10 values of $\gamma$, log spaced from $0.01 \gamma_{\max }$ to $\gamma_{\max }$
- solve all instances in 23 sec total, vs 72 sec if not sharing cache across problem instances
- (by comparison, solving one instance with CVX takes 2 minutes)


## Block splitting

- now turn to distributed setting
- approach:
(1) express graph form problems using blocks
(2) use problem transformation
(3) apply version of ADMM given earlier
(4) simplify algorithm


## Block partitioned form

- suppose $f$ and $g$ are block separable, i.e.

$$
f(y)=\sum_{i=1}^{M} f_{i}\left(y_{i}\right), \quad g(x)=\sum_{j=1}^{N} g_{j}\left(x_{j}\right)
$$

where $y_{i} \in \mathbf{R}^{m_{i}}, x_{j} \in \mathbf{R}^{n_{j}}$

- partition $A$ conformably

$$
A=\left[\begin{array}{cccc}
A_{11} & A_{12} & \cdots & A_{1 N} \\
A_{21} & A_{22} & \cdots & A_{2 N} \\
\vdots & \vdots & \ddots & \vdots \\
A_{M 1} & A_{M 2} & \cdots & A_{M N}
\end{array}\right]
$$

## Graph form problems

- graph form problem can then be expressed as

$$
\begin{array}{ll}
\operatorname{minimize} & \sum_{i=1}^{M} f_{i}\left(y_{i}\right)+\sum_{j=1}^{N} g_{j}\left(x_{j}\right) \\
\text { subject to } & y_{i}=\sum_{j=1}^{N} A_{i j} x_{j}, \quad i=1, \ldots, M
\end{array}
$$

- $M=1$ called column (feature) splitting
- $N=1$ called row (data) splitting
- goal is to solve this in a way that allows each block $A_{i j}$ to be handled by a separate process or machine


## Example: loss minimization

- row splitting corresponds to splitting by data
- column splitting corresponds to splitting by features
- loss $l$ is typically fully separable
- regularizer $r$ is often separable and sometimes block separable (e.g., group lasso or sum-of-norms regularization)


## Problem transformation

- introduce additional variables

$$
\begin{array}{lll}
\operatorname{minimize} & \sum_{i=1}^{M} f_{i}\left(y_{i}\right)+\sum_{j=1}^{N} g_{j}\left(x_{j}\right) & \\
\text { subject to } & x_{i j}=x_{j}, & i=1, \ldots, M \\
& y_{i}=\sum_{j=1}^{N} y_{i j}, & i=1, \ldots, M \\
& y_{i j}=A_{i j} x_{i j}, & i=1, \ldots, M, \quad j=1, \ldots, N
\end{array}
$$

- move some constraints into the objective

$$
\begin{array}{ll}
\operatorname{minimize} & \sum_{i=1}^{M} f_{i}\left(y_{i}\right)+\sum_{j=1}^{N} g_{j}\left(x_{j}\right)+\sum_{i=1}^{M} \sum_{j=1}^{N} I_{i j}\left(y_{i j}, x_{i j}\right) \\
\text { subject to } & x_{i j}=x_{j} \\
& y_{i}=\sum_{j=1}^{M} y_{i j}
\end{array}
$$

- now apply ADMM and simplify


## Block splitting algorithm

$$
\begin{aligned}
y_{i}^{k+1 / 2} & :=\mathbf{p r o x}_{l_{i}}\left(y_{i}^{k}-\tilde{y}_{i}^{k}\right) \\
x_{j}^{k+1 / 2} & :=\mathbf{p r o x}_{r_{j}}\left(x_{j}^{k}-\tilde{x}_{j}^{k}\right) \\
\left(y_{i j}^{k+1 / 2}, x_{i j}^{k+1 / 2}\right) & :=\prod_{i j}\left(y_{i j}^{k}+\tilde{y}_{i}^{k}, x_{j}^{k}-\tilde{x}_{i j}^{k}\right) \\
x_{j}^{k+1} & :=\operatorname{avg}\left(x_{j}^{k+1 / 2},\left\{x_{i j}^{k+1 / 2}\right\}_{i=1}^{M}\right) \\
\left(y_{i}^{k+1},\left\{y_{i j}^{k+1}\right\}_{j=1}^{N}\right) & :=\mathbf{e x c h}\left(y_{i}^{k+1 / 2},\left\{y_{i j}^{k+1 / 2}\right\}_{j=1}^{N}\right) \\
\tilde{x}_{j}^{k+1} & :=\tilde{x}_{j}^{k}+x_{j}^{k+1 / 2}-x_{j}^{k+1} \\
\tilde{y}_{i}^{k+1} & :=\tilde{y}_{i}^{k}+y_{i}^{k+1 / 2}-y_{i}^{k+1} \\
\tilde{x}_{i j}^{k+1} & :=\tilde{x}_{i j}^{k}+x_{i j}^{k+1 / 2}-x_{j}^{k+1}
\end{aligned}
$$

where $\operatorname{exch}\left(c,\left\{c_{j}\right\}_{j=1}^{N}\right)$ is given by

$$
y_{i j}^{k+1}:=c_{j}+\left(c-\sum_{j=1}^{N} c_{j}\right) /(N+1), y_{i}^{k+1}:=c-\left(c-\sum_{j=1}^{N} c_{j}\right) /(N+1)
$$

Computation: prox, $\Pi_{i j}$, dual updates

$$
\begin{aligned}
& \odot \odot \odot \odot \odot \\
& \odot \odot \odot \odot \\
& \odot \\
& \odot \\
& \odot \\
& \odot \\
& \odot
\end{aligned} \odot \odot
$$

Computation: consensus (avg)


## Computation: exchange (exch)



## Distributed lasso

- examples with dense $A_{i j} \in \mathbf{R}^{3000 \times 5000}$
- distributed solver written in C using MPI and GSL (ATLAS)
- run on Amazon EC2 cluster compute nodes
- computation times (all times in seconds)

| $M \times N$ | $4 \times 2$ | $8 \times 5$ | $8 \times 10$ |
| :--- | :--- | :--- | :--- |
| nonzero entries | 120 MM | 600 MM | 1.2 B |
| $\#$ cores | 8 | 40 | 80 |
| factorization time | 15 | 15 | 15 |
| iteration time | $0.05-0.15$ | $0.05-0.15$ | $0.05-0.15$ |
| $\#$ iterations | 90 | 230 | 490 |
| main loop time | 10 | 27 | 60 |
| total time | 28 | 50 | 80 |

## Outline

> Operator splitting
> Applications
> Block splitting

Conclusions

## Summary and conclusions

- coordinate many processors, each solving a substantial problem, to solve a very large problem
- split by data, features, or both
- yields algorithms that are easy to implement (just supply prox operators)


## Summary and conclusions

interaction between three key ingredients:
(1) proximal operator of a convex function
(2) operator splitting algorithms
(3) problem transformations

## References

- Proximal algorithms (Parikh, Boyd)
- Distributed optimization and statistical learning via the alternating direction method of multipliers (Boyd, Parikh, Chu, Peleato, Eckstein)
- Block splitting for distributed optimization (Parikh, Boyd)
available from nparikh.org or stanford.edu/~boyd

